

# Data Structures Quick Test

## Data Structures

**Q1:** Which of the following data types allows the access to the first and last element in constant time  $O(1)$ ?

- Arrays       Lists with first pointer only       Lists with pointer first and last pointer  
(Note that this question contained *removal*, instead *access to* in the previous version)

**Q2:** What is the runtime complexity to access the  $n$ th element of a list?

- $O(1)$         $O(\log(n))$         $O(n)$

**Q3:** What is the runtime complexity to access the  $n$ th element of an array?

- $O(1)$         $O(\log(n))$         $O(n)$

**Q4:** If  $f(n) \in O(g(n))$  holds then  $g(n) \in O(f(n))$  also holds.

- True       False

**Q5:** When implementing a stack using a singly-linked list, then the order of elements in the list has no impact on the runtime behaviour of the implementation.

- True       False

**Q6:** Stacks can be implemented using linked lists and arrays, whereas queues can only be implemented with linked lists, but not arrays.

- True       False

**Q7:**  $n^2 \in O(n)$  .

- True       False

**Q8:** Given a runtime function  $T(n) = 2 * n^2 + 4 * \log(n) + 5$ , which of the following statements is correct?

- $T(n) \in O(n)$         $T(n) \in O(n \log(n))$         $T(n) \in O(n^2)$         $T(n) \in O(n^2 \log(n))$

**Q9:** What is the amortized runtime behaviour of  $n$  push\_back operations if we always increase the vector size by a constant  $c$  when we resize the array?

- $O(1)$         $O(\log(n))$         $O(n)$

**Q10:** What is the amortized runtime behaviour of  $n$  push\_back operations if we always double the size of the vector every time we resize the array?

- $O(1)$         $O(\log(n))$         $O(n)$